

Layered wheels

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Dewi Sintiari, joint work with Nicolas Trotignon



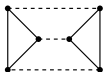
Introduction

Let G be a graph.

- H is an **induced** subgraph of G , if it can be obtained from G by deleting vertices
☞ we say that G **contains** H
- G is **H -free** if it does not contain any graph isomorphic to H
- G is **\mathcal{F} -free** if it is H -free, for every $H \in \mathcal{F}$

Truemper configurations

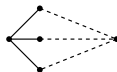
We are interested in the existence of the following graphs as induced subgraphs of the classes being studied.



prism



pyramid



theta



wheel

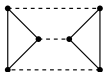
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²Conforti, Cornuéjols, Kapoor, and Vušković (2002)

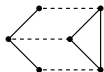
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Truemper configurations

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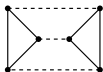
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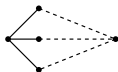
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Some classes that exclude Truemper configurations:

- **Perfect graphs** → excluding pyramid, some wheels¹
- **Even-hole-free graphs** → excluding prism, theta, some wheels²
 - ▶ even hole = a hole of even length
- **Claw-free graphs** → excluding pyramid, theta, many wheels³
 - ▶  claw
- many others...

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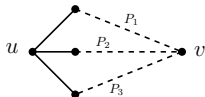
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Even-hole-free (EHF) graphs and Theta-free (TF) graphs

Recall:

- an **even hole** is a hole of even length
- a **theta** is a graph induced by three paths s.t. any two of them induce a hole.

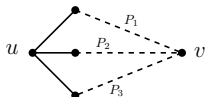


Remark. At least two of P_1, P_2, P_3 have same parity

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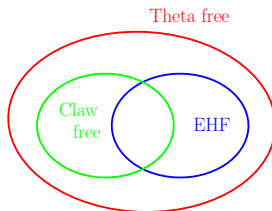
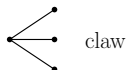
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Remark. At least two of P_1, P_2, P_3 have same parity

Observation. A **theta** always contains a **claw** and an **even hole**



EHF graphs vs TF graphs - Algorithmic/structural issues

- **Coloring**
 - ▶ NPC for TF graphs (because it is NPC for claw-free graphs ⁴)
 - ▶ open for EHF graphs
- **Max independent set**
 - ▶ open for both classes
 - ▶ **Remark.** it is polynomial for claw-free graphs ⁵
- **Decomposition theorem**
 - ▶ several decomposition theorems are known for EHF graphs ⁶
 - ▶ open for TF graphs

⁴Hoyler (1981)

⁵Minty (1980)

⁶Conforti, Cornuéjols, Kapoor, and Vušković (2002); da Silva and Vušković (2008)

Graph widths

☞ a parameter that measure how "complex" the structure of a graph is

Three notions of widths that we use:

- Treewidth tw
- Pathwidth pw
- Rankwidth rw

a well-known bound:

$$\text{For any graph } G, \quad rw(G) \leq tw(G) \leq pw(G)$$

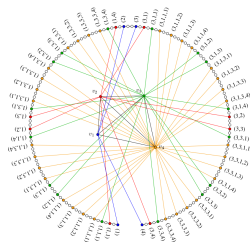
Remark.

- Observation: EHF graphs have unbounded treewidth (clique, chordal graphs are EHF)
- It is not easy to produce a (non trivial) EHF graphs of large widths
- also not easy to find a (non trivial) subclass of small widths

Some results on the widths of EHF graphs

Negative results on EHF graphs⁷

- The rankwidth of EHF graph with no diamond is unbounded



diamond

★ **Remark.** the graph contains large clique

⁷Adler, Le, Müller, Radovanović, Trotignon, and Vušković (2017)

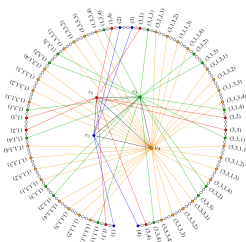
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Some results on the widths of EHF graphs

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Positive results on EHF graphs⁸

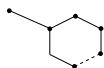
- EHF graph G with no triangle: $tw(G) \leq 5$
- EHF graph G with no cap: $tw(G) \leq 48$
- EHF graph G with no pan: $tw(G) \leq 1.5\omega(G)$

triangle



cap

= hole + short chord



pan

= hole + pendant edge

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Observation, motivation, and results

Recall: EHF graphs with no triangle have treewidth at most 5

- How about EHF graphs with no K_4 ?
 - ☞ we show that **the treewidth can be arbitrarily large**
- This also answers the following question⁹ :
is the **treewidth of EHF graph** (in general) bounded by a **function of its max clique size**?
 - ☞ no, our constructions have small clique size, but large treewidth

⁹asked by Cameron, Chaplick, and Hoàng (2018)

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Theorem 1. [DS., Trotignon (2019+)]

- For every $\ell \geq 1$, $k \geq 4$, there exists a **K_4 -free EHF-graph** of **girth at least k** and **treewidth at least ℓ** .
- idem for triangle-free TF-graph.

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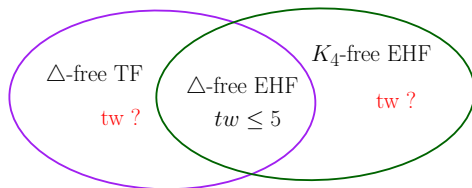
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Such graphs is what we name as **Layered wheels**

⁹asked by Cameron, Chaplick, and Hoáng (2018)

Δ -free TF graphs & K_4 -free EHF graphs

What we have and what we study...



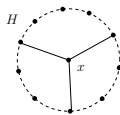
Observation:

- Δ -free TF graphs and K_4 -free EHF graphs both are (prism, pyramid, theta)-free
- Δ -free TF graphs do not contain **certain wheels**,
 K_4 -free EHF graphs do not contain **some other wheels**

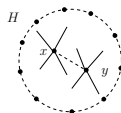
☞ Let's study the structure of **wheels** in our classes

Wheels in our classes

- **wheel** = hole H + a vertex x that has at least 3 neighbours in H
- **2-wheel** = hole H + vertices x, y that each has at least 3 neighbours in H



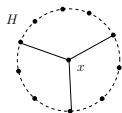
wheel



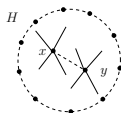
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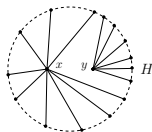
wheel



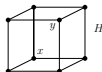
2-wheel

Structure of 2-wheels with x and y non-adjacent

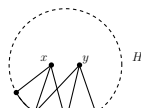
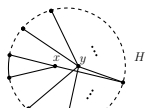
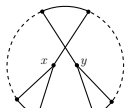
- In \triangle -free EHF graph : 2-wheels are always nested
- In \triangle -free TF graph : nested, except the cube
- In K_4 -Free EHF graph : nested, with several exceptions



nested wheel



cube



Layered wheel

Notation: Layered wheel, of ℓ layers and girth k

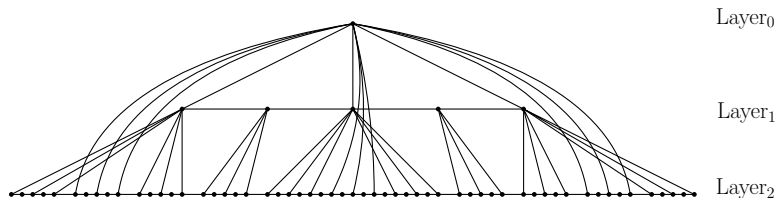


Figure: \triangle -free TF layered wheel of 2 layers with girth 4

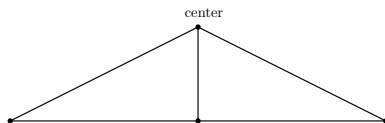
Δ -free TF layered wheel - construction

center
•

L_0

$G(\ell, k)$, with $\ell = 2$ and $k = 4$

\triangle -free TF layered wheel - construction

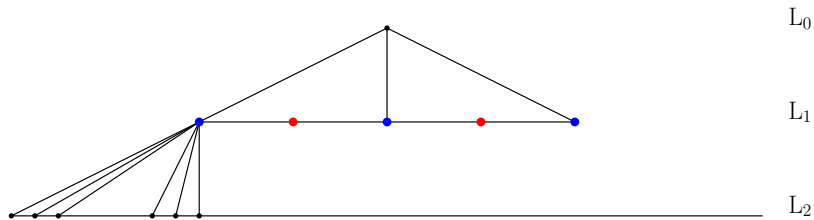


L_0

L_1

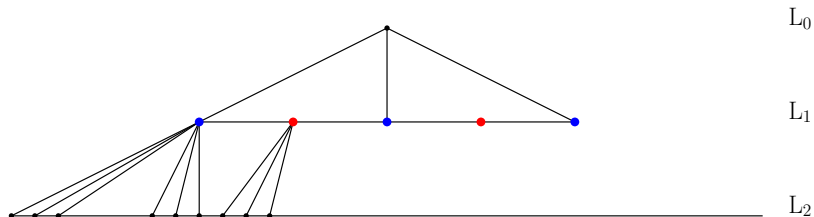
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\triangle -free TF layered wheel - construction



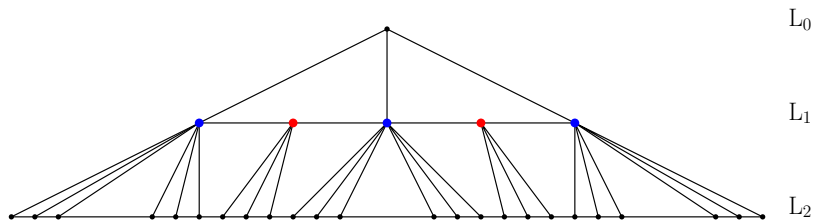
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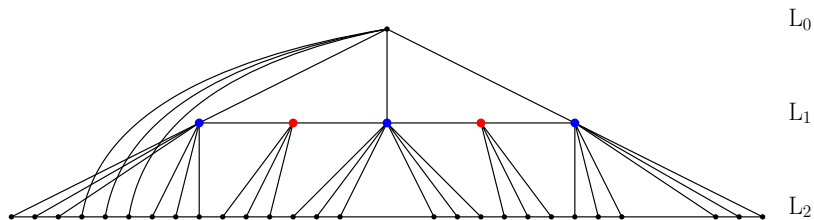
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Δ -free TF layered wheel - construction



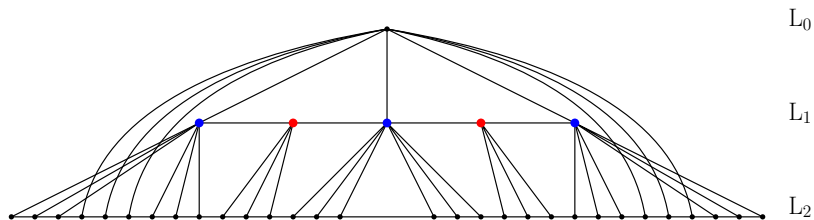
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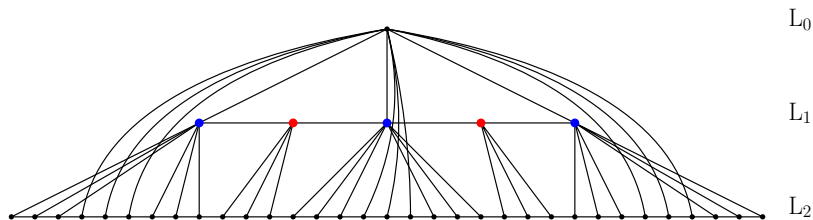
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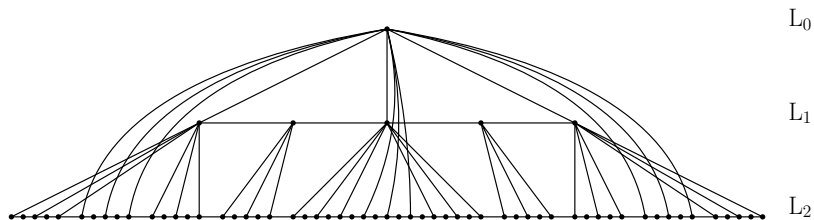
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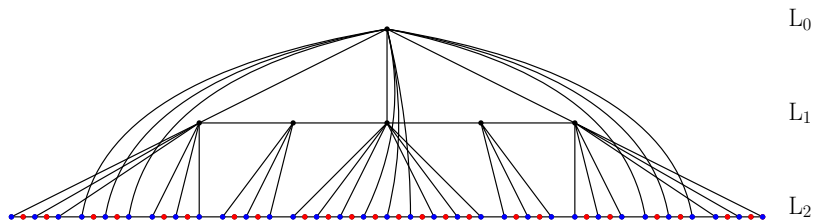
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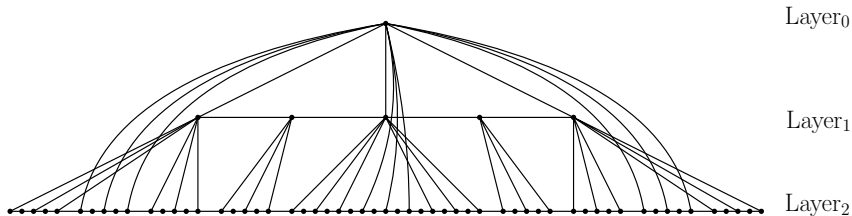
Δ -free TF layered wheel - construction



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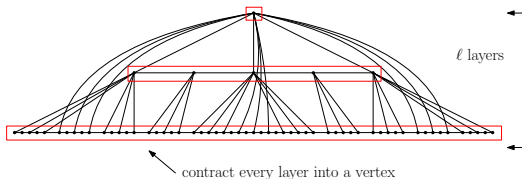
Layered wheel is theta-free

- Layered wheel is full of subdivided claw but... it contains no theta



Treewidth of layered wheel

- the treewidth of layered wheel on ℓ layers is at least ℓ



- stronger result: the rankwidth is unbounded

► **Classical theorem:** ¹⁰

Let t be integer, and \mathcal{C} be a class of graph that do not contain $K_{t,t}$ as a subgraph. Then the class has bounded tw iff it has bounded rw .

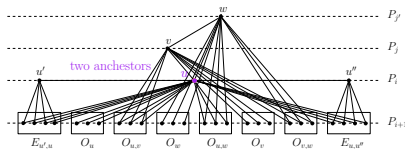
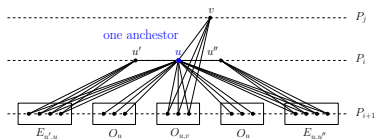
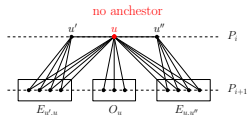


- our \triangle -free TF layered wheels have no $K_{2,3}$

¹⁰Gurski and Wanke (2000)

EHF layered wheel - construction

- The first two layers are similar as for \triangle -free TF layered wheel
- The construction is more complicated. There are three types of vertices.



- The treewidth is at least ℓ
- It is even-hole-free
- It contains no $K_{2,2}$, so rw is unbounded



$K_{3,3}$



The treewidth is "small" in some sense

Consider layered wheel on ℓ layers.

Remark. to reach treewidth ℓ , the layered wheel needs $\Omega(3^\ell)$ vertices.

Theorem 2. [DS., Trotignon (2019+)]

The treewidth of layered wheel is in $O(\log(n))$ where n is the vertex size.

Proof.

1. $n \gg 3^\ell$
2. $tw(\text{layered wheel}) \leq pw(\text{layered wheel}) \leq 2\ell$

To sum up...

Theorem 1. For every integers $\ell \geq 1$ and $k \geq 3$ there exists a graph $G_{\ell,k}$ such that:

- it is **theta-free** and it has girth at least k (so, is **Δ -free** when $k \geq 4$).
- $\ell \leq rw(G_{\ell,k}) \leq tw(G_{\ell,k}) \leq pw(G_{\ell,k}) \leq 2\ell \leq 2^\ell \leq |V(G_{\ell,k})|$.

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Theorem 2. For every integers $\ell \geq 1$ and $k \geq 3$ there exists a graph $G_{\ell,k}$ such that:

- it is **K_4 -free EHF** and every hole in the graph has length at least k .
- $\ell \leq rw(G_{\ell,k}) \leq tw(G_{\ell,k}) \leq pw(G_{\ell,k}) \leq 2\ell \leq 2^\ell \leq |V(G_{\ell,k})|$.

Logarithmic treewidth conjecture

Conjectures. [DS., Trotignon (2019+)]

- There exists a constant c such that for any Δ -free TF graph G , we have

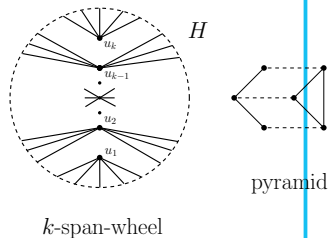
$$tw(G) \leq c \log |V(G)|.$$

- Idem for K_4 -free EHF graph.

☞ if these are true, many graph problems are poly-time solvable in both classes.

Theorem 3. [DS., Trotignon (2019+)]

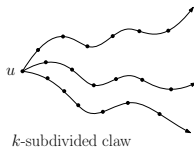
- Let k be fixed. There exists a constant c such that any (theta, triangle, k -span-wheel)-free graph G has treewidth bounded by $O(k^6)$.
- idem for (K_4 , pyramid, k -span-wheel)-free EHF graph $\rightarrow O(k^9)$



Excluding more structure

Theorem 4. [DS., Trotignon (2019+)]

Let $k \geq 1$ fixed. There exists a constant c such that any (theta, triangle, k -subdivided-claw)-free graph has treewidth bounded by $O(k^{o(1)})$.



Remark.

- The theorem is interesting regarding the max independent set problem. It is open for the class of k -subdivided-claw-free graphs and the class of theta-free graphs
- The theorem is no more true if one exclusion is forgotten.

Grid-minor-like conjecture

Grid minor theorem ¹¹

There exists a function f such that: if $tw(G) \geq f(k)$, then G contains a grid of treewidth at least k as a minor.

☞ Is there a similar theorem with **induced subgraph** instead of minor?

¹¹Robertson and Seymour (1986)

Grid-minor-like conjecture

¹²Does there exist a function f such that: every graphs with $tw \geq f(k)$ contains as an induced subgraph some graph of $tw \geq k$ that is one of the following

- a big clique
- a big complete bipartite graph
- a big grid (possibly subdivided)
- a big wall (possibly subdivided)
- a big line graph of a subdivided wall

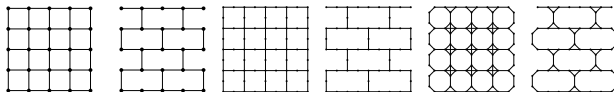


Figure: A grid, a wall, a subdivision of the former and its line graphs

¹²informally questioned by Zdeněk Dvořák

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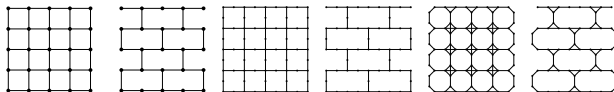


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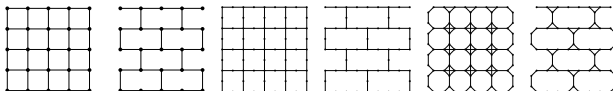


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How if some items are added to the list?

- layered wheels or some variation of them?
- long paths (P_t with $t \geq 5$)?
- a vertex of high degree (at least 4)?
- graphs containing $\geq c^f(k)$ vertices?

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- a big wall (possibly subdivided)
- a big line graph of a subdivided wall

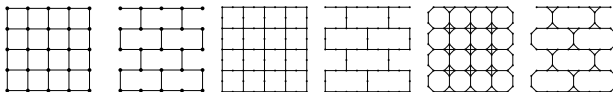


Figure: A grid, a wall, a subdivision of the former and its line graphs

no, layered wheels is a counter-example

How if some items are added to the list?

- layered wheels or some variation of them?
- long paths (P_t with $t \geq 5$)
- a vertex of high degree (at least 4)?
- graphs containing $\geq c^f(k)$ vertices?

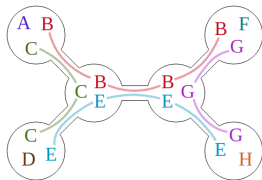
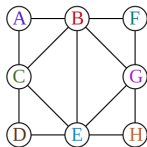


¹²informally questioned by Zdeněk Dvořák



Treewidth and pathwidth

Treewidth

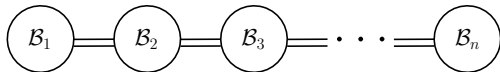


figures taken from https://commons.wikimedia.org/wiki/User:David_Eppstein/Gallery

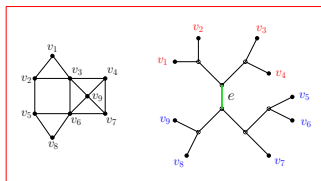
- The treewidth of G is a parameter measuring how far is G from a tree
- We create a tree decomposition whose nodes are "bags"
- The treewidth of G is the size of the largest bag minus 1 in an optimal tree decomposition

pathwidth

- "path"-version of tree decomposition



Rankwidth



$$\text{width}(e) = \text{rank} \begin{matrix} & v_1 & v_2 & v_3 & v_4 \\ v_5 & \begin{bmatrix} 0 & 1 & 0 & 0 \end{bmatrix} \\ v_6 & \begin{bmatrix} 0 & 0 & 1 & 0 \end{bmatrix} \\ v_7 & \begin{bmatrix} 0 & 0 & 0 & 1 \end{bmatrix} \\ v_8 & \begin{bmatrix} 0 & 0 & 0 & 0 \end{bmatrix} \\ v_9 & \begin{bmatrix} 0 & 0 & 1 & 1 \end{bmatrix} \end{matrix} = 3$$

example from Hlineny et. al. Width parameters beyond treewidth and their applications, The Computer Journal (51), 2008

- It is a parameter measuring the connectivity of G
- Rank decomposition is a cubic tree \mathcal{T} , with a bijection $V(G) \rightarrow \mathcal{L}(\mathcal{T})$
- the rankwidth of G is the cut-rank of the adjacency matrix of the separation in an optimal rank decomposition of G

Decomposition theorem

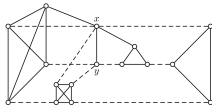
Decomposition theorem

Theorem 1. [Conforti, Cornuéjols, Kapoor, Vušković (2002); da Silva, Vušković (2008)]

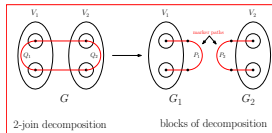
A connected ehf graph ^a is either basic or it has a 2-join or a star cutset.

^athe statement is proved for more general class, namely 4-hole-free odd-signable graph

- **Basic:** a clique, or a hole, or an extended nontrivial basic graph



- **2-join**



- **star cutset**

